# TTIC 31150/CMSC 31150 Mathematical Toolkit (Spring 2023)

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Lecture 2: Vector Space Applications and Linear Transformations

## Recap

- Fields (like  $\mathbb{R}$ ,  $\mathbb{Q}$ ,  $\mathbb{F}_p$ )
- Vector spaces (like  $\mathbb{R}^n$ ,  $\mathbb{F}_p^n$ )
- Linear dependence / independence
- Span(S)
- Basis of V
- Steinitz Exchange Principle
- Dimension of finitely-generated vector space

#### Existence of bases in general vector spaces

- Any finitely-generated vector space ( $\exists$  finite set T s.t. Span(T) = V) has a basis.
- Turns out also true for general vector spaces (even infinite-dimensional).
  - Example of such vector space? Polynomials R[X] over  $\mathbb{R}$ , or  $\mathbb{R}$  over  $\mathbb{Q}$ 
    - $f(n) = x^n$ , for  $n = 0,1,2,\cdots$
  - We define span using finite linear combination (Hamel Basis)
  - Generic vector space may not have notion of distance, closeness and convergence
- Proving it uses "Zorn's lemma" which is equivalent to axiom of choice.
- Won't get into here.

## 1 Applications of our development so far

#### 1.1 Lagrange interpolation

Lagrange interpolation is used to find the unique polynomial of degree at most n-1, taking given values at n distinct points. We can derive the formula for such a polynomial using basic linear algebra.

Recall that the space of polynomials of degree at most n-1 with real coefficients, denoted by  $\mathbb{R}^{\leq n-1}[x]$ , is a vector space. What is the dimension of this space? What would be a simple example of a basis?

• Dimension is n. Standard basis is  $\{1, x, x^2, ..., x^{n-1}\}$ .

Let  $a_1, ..., a_n \in \mathbb{R}$  be distinct. Say we want to find the unique polynomial p of degree at most n-1 satisfying  $p(a_i) = b_i \ \forall i \in [n]$ .

- Why unique?
  - ▶ If there were two, say  $p_1$ ,  $p_2$ , then  $p_1 p_2$  would have at least n roots. But a nonzero polynomial of degree at most n-1 can have at most n-1 roots.

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$$f_i(x) = \frac{g(x)}{x - a_i} = \prod_{j \neq i}^n (x - a_j),$$

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are n linearly independent polynomials in  $\mathbb{R}^{\leq n-1}[x]$ . Thus, they must form a basis for  $\mathbb{R}^{\leq n-1}[x]$  and we can write the required polynomial, say p as

$$p = \sum_{i=1}^n c_i \cdot f_i,$$

for some  $c_1, \ldots, c_n \in \mathbb{R}$ .

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for some  $c_1, \ldots, c_n \in \mathbb{R}$ . Evaluating both sides at  $a_i$  gives  $p(a_i) = b_i = c_i \cdot f_i(a_i)$ . Thus, we get

$$p(x) = \sum_{i=1}^{n} \frac{b_i}{f_i(a_i)} \cdot f_i(x).$$

other terms

evaluate to 0

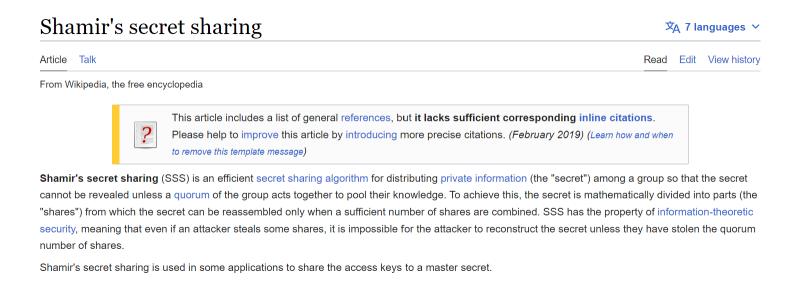
Let  $a_1, ..., a_n \in \mathbb{R}$  be distinct. Say we want to find the unique polynomial p of degree at most n-1 satisfying  $p(a_i) = b_i \ \forall i \in [n]$ .

• Argument works if replace  $\mathbb R$  with any field  $\mathbb F$  having at least n distinct points.

#### Secret Sharing

Consider the problem of sharing a secret s, which is an integer in a known range [0, M] with a group of n people, such that if any d of them get together, they are able to learn the secret message. However, if fewer than d of them are together, they do not get any information about the secret.

• E.g., password, (decryption key for) sensitive data, etc.



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- Choose a finite field  $\mathbb{F}_p$ , with  $p > \max(n, M)$ .
- Choose d-1 random values  $b_1, \ldots, b_{d-1}$  in  $\{0, ..., p-1\}$ , and let  $Q \in \mathbb{F}_p^{\leq d-1}[x]$  be the polynomial

$$Q = s + b_1 x + b_2 x^2 + \dots + b_{d-1} x^{d-1}.$$

Note that the secret is Q(0).

• For i = 1, ..., n, give person i the pair (i, Q(i)).

One direction: If any d get together, can uniquely determine Q by Lagrange interpolation, recover secret by evaluating Q at 0.

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#### Other direction:

- If d-1 get together, for any secret s', exists a consistent polynomial Q'. In fact, exactly one.
- Because Q chosen randomly from  $p^{d-1}$  polynomials consistent with secret, this means any two secrets have the same probability of producing the observed d-1 shares.

#### 3 Linear Transformations

**Definition 3.1** *Let* V *and* W *be vector spaces over the same field*  $\mathbb{F}$ . A *map*  $\varphi : V \to W$  *is called a* linear transformation *if* 

- $\varphi(v_1 + v_2) = \varphi(v_1) + \varphi(v_2) \quad \forall v_1, v_2 \in V.$
- $\varphi(c \cdot v) = c \cdot \varphi(v) \quad \forall v \in V.$

#### Example 3.2

- A matrix  $A \in \mathbb{R}^{m \times n}$  (m rows, n columns) defines a linear transformation from  $\mathbb{R}^n$  to  $\mathbb{R}^m$ . Note that we are using  $\varphi_A(v) = Av$ , where we are viewing the elements of  $\mathbb{R}^m$  and  $\mathbb{R}^n$  as column vectors.
- $\varphi$  :  $C([0,1],\mathbb{R}) \to C([0,2],\mathbb{R})$  defined by  $\varphi(f)(x) = f(x/2)$ . Recall that  $C([a,b],\mathbb{R}) = \{f : [a,b] \to \mathbb{R} \mid f \text{ is continuous}\}.$
- $\varphi$  :  $C([0,1], \mathbb{R})$  →  $C([0,1], \mathbb{R})$  defined by  $\varphi(f)(x) = f(x^2)$ .

#### Important properties

**Proposition 3.3** *Let* V, W *be vector spaces over*  $\mathbb{F}$  *and let* B *be a basis for* V. *Let*  $\alpha: B \to W$  *be an arbitrary map. Then there exists a unique linear transformation*  $\varphi: V \to W$  *satisfying*  $\varphi(v) = \alpha(v) \ \forall v \in B$ .

**Proof:** Since B is a basis, any  $u \in V$  can be written in a unique way as a sum  $\sum_{v \in B} a_v v$ , where the values  $a_v$  are in  $\mathbb{F}$  and only finitely many are nonzero. By the two properties of a linear transformation, we must then have  $\varphi(u) = \sum_{v \in B} a_v \varphi(v)$ . Since the values  $\varphi(v)$  are fixed for all  $v \in B$ , this gives the unique solution of  $\varphi(u) = \sum_{v \in B} a_v \alpha(v)$ . Moreover, this  $\varphi$  indeed satisfies the property that  $\varphi(v) = \alpha(v)$  for all  $v \in B$ .

#### Important properties

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Proposition 3.3 solidifies the connection between linear transformations and matrices. We saw that a matrix  $A \in \mathbb{F}^{m \times n}$  corresponds to a linear transformation  $\varphi_A$  from  $\mathbb{F}^n$  to  $\mathbb{F}^m$  defined as  $\varphi_A(v) = Av$ . But we can also go the other way as well. Given a linear transformation  $\varphi : \mathbb{F}^n \to \mathbb{F}^m$ , consider the standard basis  $B = \{e_1, ..., e_n\}$  for  $\mathbb{F}^n$ , where  $e_i$  has 1 in its ith coordinate and 0 in all other coordinates. By Proposition 3.3,  $\varphi$  is uniquely defined by its effect on B, and so can be represented by the matrix  $A \in \mathbb{F}^{m \times n}$  with  $\varphi(e_i)$  as its ith column.

**Definition 3.4** *Let*  $\varphi : V \to W$  *be a linear transformation. We define its* kernel *and* image *as:* 

- $\ker(\varphi) := \{v \in V \mid \varphi(v) = 0_W\}$ . [Kernel also called "nullspace"]
- $-\operatorname{im}(\varphi) = \{\varphi(v) \mid v \in V\}.$

**Proposition 3.5**  $\ker(\varphi)$  *is a subspace of V and*  $\operatorname{im}(\varphi)$  *is a subspace of W.* 

**Definition 3.6** dim(im( $\varphi$ )) is called the rank and dim(ker( $\varphi$ )) is called the nullity of  $\varphi$ .

What is rank of 
$$\varphi_A$$
 for  $A = \begin{bmatrix} 0 & 1 \\ 1 & 0 \\ 1 & 2 \end{bmatrix}$ ? Rank is 2

Nullspace just  $0_V$  since columns are independent

What is rank of 
$$\varphi_B$$
 for B=  $\begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 2 \end{bmatrix}$ ? Rank is 2 How about nullspace? All multiples of  $\begin{bmatrix} 2 \\ 1 \\ -1 \end{bmatrix}$  17

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**Proposition 3.7 (rank-nullity theorem)** *If* V *is a finite dimensional vector space and*  $\varphi : V \to W$  *is a linear transformation, then* 

$$\dim(\ker(\varphi)) + \dim(\operatorname{im}(\varphi)) = \dim(V).$$

**Proof:** Let  $n = \dim(V)$  and let  $k = \dim(\ker(\varphi))$ . Choose a basis  $v_1, ..., v_k$  for the kernel and then extend this to a basis B for V with linearly independent vectors  $v_{k+1}, ..., v_n$  (which we can always do, as we saw in the last class). We know that

$$im(\varphi) = Span(\{\varphi(v_1), ..., \varphi(v_n)\}) = Span(\{\varphi(v_{k+1}), ..., \varphi(v_n)\}).$$

So, to show that the rank is n-k, all that remains is to show that  $\varphi(v_{k+1}),...,\varphi(v_n)$  are linearly independent. This follows from the definition of linear transformation: if some linear combination of  $\varphi(v_{k+1}),...,\varphi(v_n)$  equals 0 then so does  $\varphi$  of the same linear combination of  $v_{k+1},...,v_n$ , meaning that this linear combination of  $v_{k+1},...,v_n$  lies in the kernel. This contradicts the fact that they were all linearly independent of  $v_1,...,v_k$ .